Data 101: Data Engineering Final Exam Solutions

UC Berkeley, Fall 2024 December 17, 2024

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Name of the student on your left:	
Name of the student on your right:	

Instructions

Do not open the examination until you are instructed to do so.

This exam consists of **125 points** spread over **6 questions** (including the Honor Code), and must be completed in the 170-minute time period on December 17, 2024, 3:10pm – 6:00pm unless you have pre-approved accommodations otherwise.

For multiple-choice questions, select **one choice** for circular bubble options, and select **all choices that apply** for box bubble options. In either case, please indicate your answer(s) by **fully** shading in the corresponding circle/box.

Make sure to write your SID on each page to ensure that your exam is graded.

Honor Code [2 pt]

As a member of the UC Berkeley community, I act with honesty, integrity, and respect for others. I am the person whose name is on the exam, and I completed this exam in accordance with the Honor Code.

Signature: _____

Chapter 1: What's the weather? [27 pt]

The United States Federal Emergency Management Agency (FEMA) maintains a database of federally declared disasters. We explore the two tables in this database over the next few questions.

The incidents table contains incident records	-[RECORD 1]	+
for all federally declared disasters. A sample	id	fbb556
record is shown on the right.	disaster_num	4829
Ũ	incident_type	Hurricane
Note : See the exam reference packet for a full	begin_date	2024-09-24
description of the FEMA database schema (Sec.	end_date	
1).	fips_state	45
± <i>j</i> .	declaration_title	HURRICANE HELENE

1.1. [2 pt] Suppose there are 70,000 records in **incidents** occupying a total of 5.6MB, ignoring metadata. What is the size of each record in **bytes**? Show your work.

	bytes
Solution: $\frac{5.6MB}{70,000$ records $= \frac{5.6 \times 10^6 B}{70 \times 10^3$ records $= 8$	80 bytes

1.2. [2 pt] Consider the following attribute description:

fips_state CHAR(2): FIPS two-digit numeric code used to identify the United States, the District of Columbia, US territories, outlying areas of the US and freely associated states. FIPS codes range from 00 to 99.

The attribute fips_state has been declared as CHAR(2), a fixed-length, 3-byte string type that can represent two characters. Your friend argues that a 1-byte numeric type (specifically, an integer with range -128 to 127) is more appropriate.

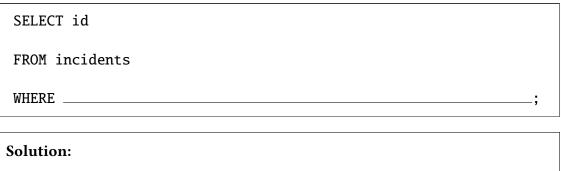
Based on the description above, would be the more appropriate type (string or 1-byte numeric) to choose for the fips_state attribute? Fill in the appropriate bubble and justify your choice in **one sentence**.

A \bigcirc string / \bigcirc 1-byte numeric data type is more appropriate for fips_state.
Reason:

Solution: A string data type is more space, but could be appropriate because it communicates that a FIPS state code is not an attribute which you can add or subtract like an integer. Otherwise, a (1-byte) integer may be more appropriate simply because it is a more efficient storage mechanism.

Note: Please read the full database schema in the exam reference packet (Sec. 1) before continuing. We have included relevant attribute descriptions below:

- disaster_num INTEGER: Sequentially assigned number used to designate an event or incident declared as a disaster.
- incident_type TEXT: The primary or official type of incident, e.g., fire, flood, etc.
- **begin_date** DATE: Date the incident began.
- **1.3.** You are asked to design dbt-style tests. For each statement below, write a query that returns an empty table if the statement is true for all records; otherwise, return the ids of records for which the statement is **false**. You may not need all blanks.
 - i. [2 pt] *All incidents have a begin date.* Write a query that returns the IDs of incidents that have a NULL begin date, or an empty table if no records match.



SELECT id FROM incidents WHERE begin_date IS NULL;

ii. [6 pt] All incidents with the same disaster number have the same begin date.Write a query that returns the IDs of incidents that have the same disaster number but *different* begin dates, or an empty table if no records match.

Note: Equality operators (=, !=) work with date type attributes.

WITH min_dates AS (
SELECT	
OVER () AS val
FROM incidents) SELECT id FROM incidents JOIN min_dates	
ON	
WHERE	;

Solution:

```
WITH min_dates AS (
   SELECT DISTINCT disaster_num, -- DISTINCT optional but more efficient
   MIN(begin_date)
        OVER (PARTITION BY disaster_num) AS val
   FROM incidents
)
SELECT id
FROM incidents
JOIN min_dates ON min_dates.disaster_num = incidents.disaster_num
WHERE begin_date != min_dates.val;
```

1.4. [2 pt] Consider the functional dependency $X \rightarrow Y$. Specify X and Y such that this functional dependency is equivalent to the statement in Question 1.3(ii).

$X: \bigcirc \mathbf{A}$.	id	$Y: \bigcirc \mathbf{A}$.	id
○ B.	disaster_num	○ B.	disaster_num
○C.	begin_date	⊖ C .	begin_date

1.5. [2 pt] According to the National Oceanic and Atmospheric Administration (NOAA), incidents like Hurricanes and Tropical Storms are given short, distinctive names to avoid confusion and streamline communications. What is the Levenshtein distance between the strings *Irene* and *Helene*?

Levenshtein distance:	3

1.6. [4 pt] The incident type (incident_type attribute) is the officially designated type of a given incident, such as fire, flood, and so on. Find the number of unique disasters (i.e., disaster numbers) per incident type, sorted by most unique disasters first, *[exam clarification]* meaning sorted by the highest number of unique disasters with that incident type.

incident_type	count
Fire	1648
Severe Storm	1097
	•••

Sample output on right. You may not need all blanks.

SELECT	
FROM incidents	
	;

Solution:
<pre>SELECT incident_type, COUNT(DISTINCT disaster_num) FROM incidents</pre>
GROUP BY incident_type ORDER BY COUNT(DISTINCT disaster_num) DESC;

1.7. [2 pt] There are 26 incident types (Fire, Flood, Hurricane, Tropical Storm, etc.) Assuming that all records have non-NULL incident types, what is theoretically the minimum number of bits needed to encode (i.e., represent) the incident_type attribute? Ignore any practical limits of SQL.

Note: See the exam reference packet for a base-2 logarithm lookup table (Sec. 10.1).

 \bigcirc 4 bits \bigcirc 5 bits \bigcirc 8 bits \bigcirc 26 bits \bigcirc None of these

1.8.	[5 pt] Suppose that all incidents have an begin date, and all
	incidents with the same disaster number have the same begin
	date (see Question 1.3). However, not all incidents have an
:	incident end date .

You decide to **impute missing end dates using the average duration (in days)** of disasters by incident type, which has been already computed for you in a view, **avg_durations**.

incident_type	avg
Fire	15
Flood	13
Hurricane	13

avg_durations

٦

Write a query that uses avg_durations to create a column, impute_end_date, that takes the end date (if it exists) or imputes the end date as begin date plus average duration of incident per type:

disaster_num	incident_type	begin_date	end_date	impute_end_date
4829	Hurricane	2024-09-24		2024-10-06
4837	Tropical Storm	2024-09-15	2024-09-19	2024-09-19
5534	5534 Fire			2024-09-15
5534	Fire	2024-09-01		2024-09-15
4037	Hurricane	2011-08-24	2011-08-30	2011-08-30

Fill in the blanks such that the below query outputs records like those above.

Notes: You may not need all blanks. See the **exam reference packet** for an example of **CASE (Sec. 6.2)**. The addition (+) operator supports date types: 10 plus 2024-10-01 is 2024-10-11.

ELECT disaster_num, incidents.incident_type, begin_date, end_date
(CASE
END) AS impute_end_date
ROM incidents
;

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Solution:
<pre>SELECT disaster_num, incidents.incident_type, begin_date, end_date, (CASE WHEN end_date IS NULL THEN (begin_date + avg_durations.avg - 1)</pre>
ELSE end_date
END
) AS impute_end_date
FROM incidents
JOIN avg_durations
<pre>ON avg_durations.incident_type = incidents.incident_type;</pre>

Chapter 2: Outliers Gonna Outlie [24 pt]

We continue our exploration of the FEMA dataset with the grants		
table, which contains financial assistance information about disas- ters. Consider the following attribute description:		value
		1
	5	160.8
ia_approved_num INTEGER: The number of disaster	10	262.6
assistance applications that were approved for Individ-	50	1984
ual Assistance (IA).	90	28022.6
Demonstile statistics for this statilests and shows to the wight	95	77026.8
Percentile statistics for this attribute are shown to the right. You would like to detect outliers on this attribute.	100	774691

2.1. Use the percentile table to determine each summary statistic below. If impossible, fill in the "Cannot be determined" bubble and leave the line blank.

i. [1 pt] Mean	N/A	0	Cannot be determined
ii. [1 pt] Median	1984	0	Cannot be determined
iii. [1 pt] Minimum	1	\bigcirc	Cannot be determined

2.2. [3 pt] Suppose that the Median Absolute Deviation (*MAD*) of ia_approved_num is *MAD* = 1690. The Hampel X84 outlier defines outliers as values that are $2k \cdot MAD = 5011.188$ from the median, where k = 1.4826 and *MAD* is the Median Absolute Deviation. With this method, which of the below ia_approved_num values are outliers?

A. 13	🗌 D. 5956	
🗌 B. 196	🗌 E. 6750	G. 49267
C. 1984	F. 8084	H. 108986
		☐ I. Cannot be determined

2.3. Suppose we instead would like to apply a 10% winsorization of ia_approved_num (i.e., 10% tails on each side). For each of the below values, what would be its corresponding transformed value in the resulting 10% winsorized array? If the correct value is not listed, fill in the "Another value" bubble.

i. [2 pt] 13	○ 13○ 160.8	<u> </u>	<mark>262.6</mark> Another value
ii. [2 pt] 5956	19845956	00	28022.6 Another value
iii. [2 pt] 49267	4926728022.6	<u> </u>	77026.8 Another value

Finally, we consider the relationship between the two tables in the FEMA dataset:

- The incidents table contains incident records for all federally declared disasters.
- The grants table contains financial assistance information about disasters.

Both tables have the attribute disaster_num, which is a sequentially assigned number ("disaster number") used to designate an event or incident declared as a **disaster**.

2.4. You use psql to explore the relationships between entities: incidents, grants, and disasters).

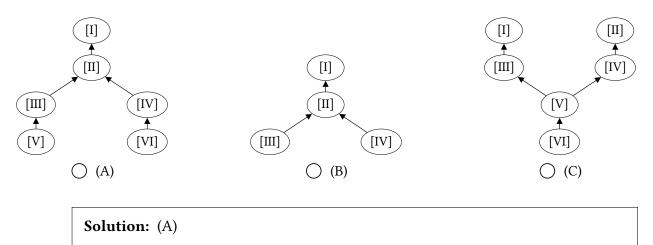
fema=#	SELECT COUNT(*),	<pre>fema=# SELECT COUNT(*),</pre>
fema-#	COUNT(DISTINCT disaster_num)	<pre>fema-# COUNT(DISTINCT g.disaster_num),</pre>
fema-#	FROM grants;	<pre>fema-# COUNT(DISTINCT i.disaster_num)</pre>
count	count	fema-# FROM incidents AS i
	-+	fema-# JOIN grants AS g
3500	3500	fema-# ON g.disaster_num
		<pre>fema-# = i.disaster_num;</pre>
fema=#	SELECT COUNT(*),	count count count
fema-#	COUNT(disaster_num),	
fema-#	COUNT(DISTINCT disaster_num)	50000 3500 3500
fema-#	FROM incidents;	
count	count count	
	-+	
70000	70000 5000	

Given the above output, mark each of the below statements as True or False.

i. [1 pt] Each grant is associated with one disaster.	<mark>○ True</mark> ○ False
ii. [1 pt] Each disaster has an associated record in the grants table.	○ True ○ False
iii. [1 pt] One disaster number can be associated with many incidents.	○ True ○ False
iv. [1 pt] The disaster_num attribute can be a pri- mary key for the grants table.	○ True ○ False
. Suppose we write the query to the right. A query optimizer then produces an optimized query plan according to SQL query semantics and projec- tion pushdown . What is the resulting optimized query plan?	<pre>SELECT incidents.id, grants.id, incident_type, ia_approved_num FROM incidents JOIN grants ON incidents.disaster_num = grants.disaster_num;</pre>

2.5

i. [2 pt] First, select the tree shape that **most closely** resembles an optimized plan (for this query) that leverages projection pushdown:



- ii. [6 pt] For the query plan you selected above, define each node by filling in the blanks:
 - For relational operators, write in the relational operator and its subscript.
 - For scans, write the table name.
 - Depending on the tree shape you selected, you may not need all blanks below. Write *N*/*A* on the blanks that you do not need.

Chapter 3: I'm Craving Mongolian Food! [19 pt]

For this question, we'll be using the Yelp dataset, looking just at the user and business entities. We'll start by using MongoDB to find some prolific Yelp users who might have good reviews or recommendations. When completing the MongoDB queries, use MQL syntax, and assume that all documents have been loaded into the db.user collection.

Consider the following properties of a user document. For this dataset we will assume all users have the same possible fields, but some fields may be null if the data does not apply to that user. You can review more about the dataset if you wish in the **exam reference packet (Sec. 2)**.

A User Document

{

}

```
"_id": ObjectId("6758dca1091ea6eafadeb758"),
"user_id": "q_QQ5kBBwlCcbL1s4NVK3g",
"name": "Jane",
"review_count": 1220,
"yelping_since": "2005-03-14 20:26:35",
"useful": 15038,
"funny": 10030,
"cool": 11291,
"elite": [2011 2012, 2013, 2014], // or null
"friends": ["friend_id_1", "..."], // or null
"fans": 1357,
"average_stars": 3.85
```

3.1. Find the total number of elite Yelp users for each year.

In order to complete this query, we will need to group users by the years in which a user is "elite". Yelp has some criteria by which a user is considered "elite". A user is an "elite" user for 2023 if the year "2023" appears in the elite field. The value "eliteCount" will be the total number of users who are elite in that year. (We will ignore any users who are not elite in any year, which are represented by a NULL value for the elite field.)

The document returned by this query should have the following structure:

```
Count of Elite Users by Year Result
[
    { "_id": 2023, "eliteCount": xxx },
    { "_id": 2022, "eliteCount": yyy },
    // "_id" is the year, one entry per year.
    // and so on...
]
```

Fill in the blanks below so that the query runs correctly.

Count of Elite Users by Year Query

```
db.user.___(i)___([
  {
    ___(ii)___: {
            $elite: {
            $exists: true,
            $ne: []
      }
    }
  },
  { ___(iii)___: "$elite" },
  {
    $group: {
      _id: {
        $toInt: ___(iv)___ },
        eliteCount: { ____(v)___: 1
      }
    }
  }
])
```

Solution:

```
db.users.aggregate([
  {
    $match: {
      elite: {
        $exists: true,
        $ne: []
      }
    }
  },
  { $unwind: "$elite" },
  {
    $group: {
      _id: { $toInt: "$elite" },
      eliteCount: { $sum: 1 }
    }
  }
])
```

i. [1 pt] () find () aggregate () groupBy () collect

ii. [2	[2 pt]	
	Solution: \$match	
iii. [<pre>[1 pt]</pre>	
iv. [2	[2 pt]	
	Solution: "\$elite"	
v. [[1 pt] 🔿 \$count 🔿 \$total 🔿 \$max 🔿 \$add 🔿 \$sum	

3.2. For this question, consider the business entity.

The complete Yelp dataset includes millions of businesses. However, we are interested in exploring alternative ways of storing data about businesses that might be more efficient. Consider the following properties of one business:

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A Business Document

```
{ "_id": ObjectId("6758dc70091ea6eafadc26c3"),
 "business_id": "6iYb2HFDywm3zjuRg0shjw",
 "name": "Bear's Lair",
 "city": "Berkeley",
 "state": "CA",
 "postal_code": "94720",
 "stars": 3.3,
 "review_count": 102,
 "is_open": 1,
 "attributes": {
    "RestaurantsTableService": "True",
    "BusinessAcceptsCreditCards": "True",
    . . .
    "RestaurantsDelivery": "GrubHub"
  },
 "categories": ["Gastropubs", "Food", ...],
 "hours": {
    "Monday": "11:00-23:00",
      "Tuesday": "11:00-23:00",
      "Wednesday": "11:00-23:00",
      "Thursday": "11:00-23:00",
      "Friday": "11:00-23:00",
      "Saturday": "11:00-23:00",
      "Sunday": "11:00-23:00"
    }
}
```

- i. [3 pt] Which fields for a given business would need to be adapted to fit into a purely relational model? i.e., without modification, these fields would not represent atomic data. (Select all that apply)
 - A. hours
 - 🗌 B. stars
 - C. review_count
 - □ D. attributes

E. is_open

- ☐ F. postal_code ☐ G. categories
- tributes

ii. [3 pt] Which of these fields represent data which is denormalized? (Select all that apply)

- \Box A. hours
- □ B. stars
 - D. Stars
- C. review_count
- □ D. attributes

- 🗌 E. is_open
- □ F. postal_code
- □ G. categories
- **3.3.** [6 pt] We want to model the hours field as a relational table, called business_hours. It will store the hours of each business, where **each row represents the hours a business** is open for a one day of the week, with columns for the opening and closing times stored separately.

Fill in the following table with the column name, data type, and whether the column is a primary key, foreign key, or neither. You should fill all the rows. We have provided some values, but you should expect to fill all blanks.

- PK means "primary key", FK means "foreign key"
- Remember each row needs a primary key
- You can assume that a **businesses** table contains and ID for each business in the typical format (e.g., "6iYb2HFDywm3zjuRg0shjw").
- There are multiple valid responses for column names and types.

Column Name	Column Type	Constraints
	integer	\bigcirc PK \bigcirc FK \bigcirc N/A
	 ○ BOOLEAN ○ INTEGER ○ TEXT ○ TIMESTAMP 	\bigcirc PK \bigcirc FK \bigcirc N/A
day_of_week	○ BOOLEAN ○ INTEGER○ TEXT ○ TIMESTAMP	\bigcirc PK \bigcirc FK \bigcirc N /A
	OBOOLEANOINTEGEROTEXTOTIMESTAMP	\bigcirc PK \bigcirc FK \bigcirc N /A
	OBOOLEANINTEGEROTEXTTIMESTAMP	\bigcirc PK \bigcirc FK \bigcirc N /A

Chapter 4: Oh! That's Lots of Books! [28 pt]

We're headed back to the library! Since the midterm, our library has gotten very busy. Due to all the traffic, the schema for our library has evolved a bit:

books (id, title, author, isbn, publication_year) book_copies (id, book_id, location_id, format, acquisition_date, status); users (id, first_name, last_name, email, phone_number, joined_date, is_ya) locations (id, name, address, phone_number) checkouts (id, user_id, book_copy_id, location_id, checkout_date, due_date) book_returns (id, checkout_id, location_id, return_date)

In this schema, each user checks out a particular book_copy which creates an entry in the checkouts table. When a user returns a book, an entry is created in the book_returns table. Note: The full library schema is included in the exam reference packet (Sec. 3).

4.1. [8 pt] Checked out books revisited.... Write a query to return the checkout_id, user's first_name, book title, book_copy_id, and due_date for books which have not yet been returned. (You may not need all blanks.)

The resulting table should look like this:

<pre>checkout_id </pre>		title +	book_copy_id	
4	Jonathan	Basic Chemistry Basic Chemistry	2802	2020-06-13 2021-03-09

SELECT c.id as checkout_id,

FROM checkouts c
JOIN users u ON c.user_id = u.id
JOIN book_copies bc
ON c.book_copy_id = bc.id
JOIN books b ON bc.book_id = b.id
WHERE NOT EXISTS (
__________);

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```
Solution:
```

```
SELECT
c.id as checkout_id,
    u.first_name, b.title, c.book_copy_id, c.due_date
    FROM checkouts c
    JOIN users u ON c.user_id = u.id
    JOIN book_copies bc ON c.book_copy_id = bc.id
    JOIN books b ON bc.book_id = b.id
    WHERE NOT EXISTS (
        SELECT id FROM book_returns r
        WHERE r.checkout_id = c.id
    );
```

4.2. [6 pt] You are tasked with **finding the most popular books**, when they are checked, and at which location, which will add the library in the purchasing of new copies of books.

We want to create a view which will aggregate all books (book_id), the location (location_id), year, month (of the checkout date) as well as the total copies_checked_out in that month. The view will have a schema that looks like this:

			•	copies_checked_out
			1	8
28	3	2020	2	9

Fill in the blanks so that checkouts_olap will match this description. Note: The SQL function EXTRACT returns just the year, month, day, etc. part from a date datatype, e.g., EXTRACT(YEAR FROM '1995-09-05') returns 1995.

CREATE OR REPLACE VIEW checkouts_olap AS SELECT

```
EXTRACT(YEAR FROM c.checkout_date) as year,
EXTRACT(MONTH FROM c.checkout_date) as month,
```

_____ as copies_checked_out

```
FROM checkouts c
JOIN book_copies bc ON c.book_copy_id = bc.id
JOIN books b ON bc.book_id = b.id
GROUP BY
```

```
EXTRACT(YEAR FROM c.checkout_date) as year,
EXTRACT(MONTH FROM c.checkout_date) as month
ORDER BY book_id DESC;
```

```
Solution:
```

```
CREATE OR REPLACE VIEW checkouts_olap AS
SELECT
    b.id as book_id,
    c.location_id,
    EXTRACT(YEAR FROM c.checkout_date) as year,
    EXTRACT(MONTH FROM c.checkout_date) as month,
    COUNT(*) as copies_checked_out
FROM checkouts c
JOIN book_copies bc ON c.book_copy_id = bc.id
JOIN books b ON bc.book_id = b.id
GROUP BY
    b.id,
    c.location_id,
    EXTRACT(YEAR FROM c.checkout_date),
    EXTRACT(MONTH FROM c.checkout_date)
ORDER BY book_id DESC;
```

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Note: As part of the course content, we did not cover CUBE() and ROLLUP(), and so we do not expect you to know how to use these PostgreSQL functions right off-the-bat. However, one of the learning goals of this course is to read SQL documentation to use new functions. **The SQL documentation for CUBE() and ROLLUP() is in the reference packet (Sec. 4)** and is sufficient for you to answer the below question.

4.3. [5 pt] Now we'd like to use the checkouts_olap view to roll up data, exploring all possible combinations of book_id, location_id and year, along with the sum of all checkouts in that group.

The resulting table will look like this:

	location_id year	
28	3 2020	
28	 3 All	300
28	All All	300
A11	All All	300
A11	All 2020	113
28	All 2020	113
A11	3 All	300
A11	3 2020	113

Fill in the blanks of the query below. You can receive full credit for this query even if the previous query is not correct.

Note: The function COALESCE(x, y) replaces a NULL values of x with the value y, so that it is more clear what the resulting group represents.

<pre>SELECT COALESCE(book_id::text, 'All') as book_id, COALESCE(location_id::text, 'All') as locati COALESCE(year::text, 'All') as year,</pre>	.on_id,
	AS total_checkouts
FROM	
GROUP BY	;
Solution:	
SELECT	

rows

```
COALESCE(book_id::text, 'All') as book_id,
COALESCE(location_id::text, 'All') as location_id,
COALESCE(year::text, 'All') as year,
SUM(copies_checked_out) as total_checkouts
FROM checkouts_olap
GROUP BY CUBE(book_id, location_id, year);
```

4.4. [2 pt] Assuming checkouts_olap contains 4 unique book_ids, 3 unique location_ids, and 4 unique years, how many total rows will be in the resulting table from Question 4.3? Briefly, show how you calculated this value.

Solution: When using CUBE we have the number of rows plus one more for the cumulative total. So for 4 book_ids, 3 location_ids, and 4 years, we have $5 \times 4 \times 5 = 100$ total rows.

4.5. Starting with the data in checkouts_olap, which of the types of roll up or drill down operations would be possible?

i.	[1 pt] Roll up total checkouts for a given book across all library locations.	🔿 Yes	🔿 No
ii.	[1 pt] Roll up total checkouts by the user's joined_date after joining with the users table.	⊖ Yes	🔿 No
iii.	[1 pt] Drill down into the day each book copy was checked out.	🔿 Yes	🔿 No

- **4.6.** When a user checks out a particular **book_copy**, a transaction T_i is executed with the following ordered actions:
 - 1. $R_i(B_c)$: Read from **book_copies** (to check status available).
 - 2. $W_i(B_c)$: Write to **book_copies** (to update status to checked out).
 - 3. $W_i(C)$: Write to checkouts (to insert a new entry).

Suppose that two users simultaneously check out the same book_copy, resulting in concurrent transactions T_1 and T_2 . The DBMS generates the following transaction schedule:

ALL 202	4		Final E	XAM		SID:
	1	2	3	4	5	6
T_1	$R_1(B_c)$	$W_1(B_c)$			$W_1(C)$	
T_2			$R_2(B_c)$	$W_2(B_c)$		$W_2(C)$

- ii. [2 pt] Does this schedule provide isolation be-tween T_1 and T_2 ?
- O Yes

🔘 No

O Cannot be determined

Chapter 5: Teenagers! At The Library [12 pt]

One of the librarians wants to understand how many books "young adults" (i.e., teenagers) are checking out at the library in the month of December. They construct a **spreadsheet**:

	Α	В	С	D	E	F
1	id	date	user_id	num_out	is_ya	ya_tot
2	001	2024-12-01	101	1	TRUE	1
3	002	2024-12-01	980	2	TRUE	3
4	003	2024-12-02	103	1	FALSE	3
5	004	2024-12-14	980	2	TRUE	5

- The topmost row and leftmost column are the column and row addresses, respectively, of the spreadsheet (e.g., cell *B*4 has the value 2024-12-02).
- Each row refers to the **total** number of checkouts num_out made by a user with ID user_id on a specific date date in December 2024. The is_ya flag is TRUE only if the user is a young adult. id is a primary key.
- Rows are sorted by earliest date first; ties are broken by user_id.

The librarian has already pre-computed columns A–E but would like your help with computing the **young adult total** (column *F*, *italics*): a cumulative sum of the number of checkouts (column *D*) made by young adults (column *E*) up through the current row.

5.1. [2 pt] Use a the SUMIF formula to compute Column *F*, rows 2 onwards. From the Google Sheets documentation, SUMIF(range, criterion, [sum_range]) returns a conditional sum across a range. Note: See the exam reference packet for more detailed documentation (Sec. 5).

Specify the correct formula to insert into cell *F*2 such that, when dragged down (i.e., filled down), cells *F*2 downwards will achieve the desired behavior.

- A. =SUMIF(D\$2:D2,E\$2:E2)
 B. =SUMIF(E\$2:E2,D\$2:D2)
 C. =SUMIF(D\$2:D2, TRUE, E\$2:E2)
 D. =SUMIF(E\$2:E2, TRUE, D\$2:D2)
 - \bigcirc E. None of the above
- **5.2.** [6 pt] Consider how this view would be implemented in different data models. In particular, spreadsheet models support both row addressing and column addressing. Which of the following models support each type of addressing? Select all that apply.

i. PostgreSQL tables	□ Row addressing	Column addressing] Neither
ii. pandas DataFrames	Row addressing	Column addressing	☐ Neither
iii. 2-D tensors	Row addressing	Column addressing	☐ Neither

You consider instead computing the "young adult total" column using the original PostgreSQL library database schema.

You first construct a view named dec_checkouts that computes the attributes id, date, user_id, num_out, and is_ya for checkouts in December 2024. The view dec_checkouts is the SQL equivalent of columns A-E in the librarian's spreadsheet.

5.3. [4 pt] Your friend who has not taken Data 101	SELECT id, date, user_id,
reads the PostgreSQL documentation and writes	num_out, is_ya,
the query to the right:	SUM(num_out) OVER (
	PARTITION BY is_ya
This query computes the librarian's requested	ORDER BY id ASC) AS ya_tot
ya_tot column with a window function.	FROM dec_checkouts
yu_coc column while a window function.	ORDER BY date ASC;

Your friend's query produces the following output:

id	date		user_id				-	-	
	024-12-01							 	1
002 2	024-12-01	I	980	I	2	I	t		3
003 2	024-12-02	I	103	I	1	l	f		1
004 2	024-12-14	I	980	I	2	l	t	I	5

The output looks **almost** correct but has **an incorrect value of 1** for ya_tot in record ID 003 (see spreadsheet cell *F*4, which has value 3).

In no more than **two sentences**, explain to your friend why their query produces this incorrect value.

Solution: Window functions can only compute (aggregate) functions over their specified partition. Your friend's code is effectively computing *two* cumulative sums, for both young adults and non-young adults.

Chapter 6: A Not-So-Random Sampling of Topics [13 pt]

For each of the below multiple-choice questions, select **one choice** if circular bubble options, and select **all choices that apply** if box bubble options. In either case, please indicate your answer(s) by **fully** shading in the corresponding box/circle.

6.1.		e	supports JOINing data across different documents clarification].	0 0	True False		
6.2.	[1 pt]	MongoDB o	does not support indexing data.	0	True	\bigcirc	False
6.3.	[1 pt]	MongoDB i	is an example of a document store.	\sim	True False		
6.4.	[1 pt]	PostgreSQI	cannot support semi-structured data like JSON.	<mark>0</mark> 0	True False		

- **6.5.** [2 pt] Suppose we would like to distribute our data across multiple servers. An RDBMS like PostgreSQL is preferred over a NoSQL system like MongoDB when:
 - ☐ A. Your data schema is normalized.
 - B. Your data requirements and structure are constantly changing.
 - C. You have a large amount of sparse data.

D. It is imperative that all data follow very strict constraints.

6.6.	[1 pt] It is possible to shard (partition) any database in a horizontal model, regardless of the structure of the data.	\sim	True False		
6.7.	[1 pt] It is possible to shard (partition) any database in a vertical model, regardless of the structure of the data.	0	True	0	False
6.8.	[1 pt] Memory and disk refer to the same physical storage element on a computer.	0	True	0	False
6.9.	[2 pt] In a highly parallel MapReduce system, each mapper process	\bigcirc	True	\bigcirc	False

- gains performance by read/writing data from other mapper processes.
- **6.10.** [2 pt] Which ACID property is best illustrated by this situation: In a highly concurrent system, many users read from and write to the same bank account statements simultaneously.

○ Atomicity ○ Consistency ○ Isolation ○ Durability

Chapter 7: Congratulations! [0 pt]

Congratulations! You have completed this exam.

- Make sure that you have written your Student ID number on every other page of the exam. You may lose points on pages where you have not done so.
- Also ensure that you have signed the Honor Code on the cover page of the exam for 1 point.
- If more than 10 minutes remain in the exam period, you may hand in the exam **and** the reference packet and leave.
- If \leq 10 minutes remain, please sit quietly until the exam concludes.

[Optional, 0 pts] Use this page to draw your favorite Data 101 moment!